COMFAST: A Comparative Framework for Analysis of Scheduling Techniques in Multi cores

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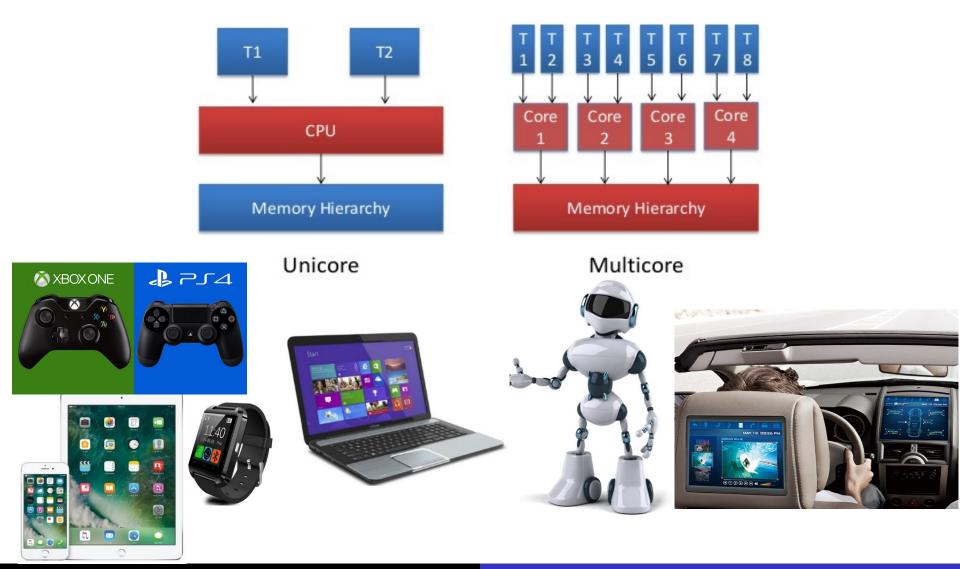
Introduction and Motivation

Proposed Methodology: COMFAST

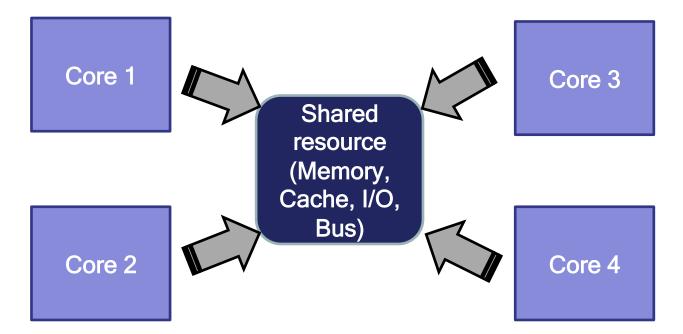
Results

Conclusions

Multi-core Architectures



Resource Contention Problem



- Resource contention can lead to the accumulation of delays, which may lead to
 - Degraded performance
 - Tasks missing their deadlines

Scheduling Algorithms

- Used to allocate tasks to cores, ensuring that the time constraints of all tasks are met
 - Load Balancing [Merkel, 2010]
 - Improve system's performance by migrating tasks from busy cores to idle cores, at the cost of additional migration time

Reasonable Allocation Techniques [Huang, 2016]

- First Fit (first core with available capacity),
- Best Fit (core with maximum available capacity)
- Worst Fit (core that has the minimum capacity available)

Task Aware Scheduling [Chiang, 2016]

 Allocate tasks based on their resource requirements to ensure optimal utilization for resources

Evaluation of Scheduling Algorithms

□ Worst Case Response Time (WCRT)

Sum of its execution time and the delays that are accumulated due to preemptions from higher priority tasks

Most of the existing analysis approaches make their preferable assumptions for the tasks and system resources

Fail to compare different scheduling techniques on a common and fair ground

Proposed Solution: COMFAST

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 System Model
 A generic Analysis Approach

COMFAST: System Model

$\Box m$ cores

Task graph: A set of n tasks

$$\Box G_{i} = \{T_{1}, T_{2}, \dots T_{n}\}$$

T: task

R is calculated using τ^{WCET} and the surplus time added due to resource contention

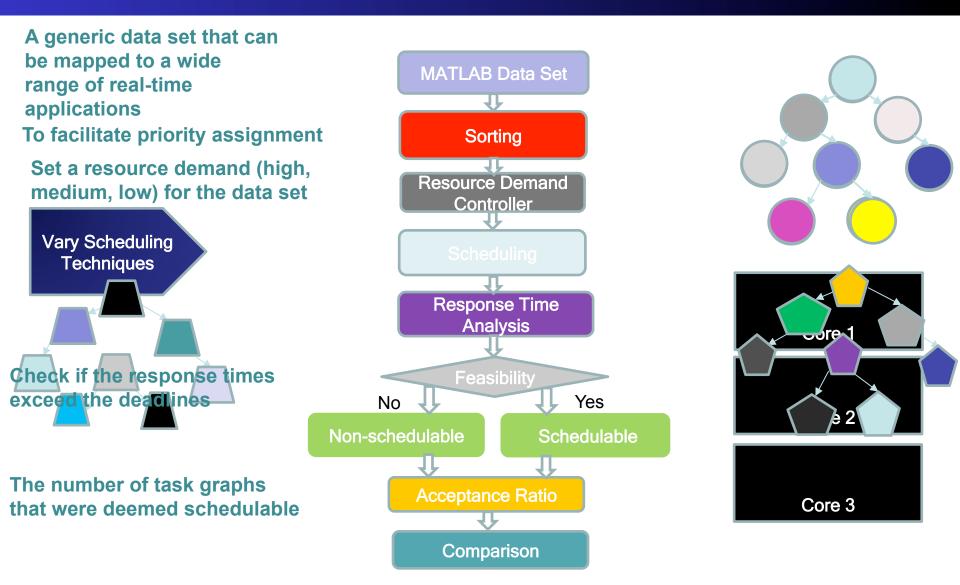
$$\tau^d = \tau^N * \tau^M$$

D_i	Task deadline	
τ^{BCET}	Best case execution time	
τ^{WCET}	Worst case execution time	
R_i	Response time	
τ^{minS}	Minimum start time	Define
τ^{maxS}	Maximum start time	Define
τ^{minF}	Minimum finish time	dependencies
τ^{maxF}	Maximum finish time	between tasks
τ^N	Maximum number of resource accesses	
τ^M	Maximum access duration	
$ au^d$	Total resource demand	Deadline
τ^{pri}	Task priority	monotonic
τ^p	Processor	priority

COMFAST: System Model

- □ A task graph is represented by its schedulability (G_S)
 - □Schedulable (G_S = 1): If all its constituent tasks meet their deadlines
 - □Non-schedulable ($G_S = 0$): If any one of its tasks is unable to meet its deadline
- Data set is characterized by acceptance ratio (σ)
 Ratio of the number of schedulable task graphs to the total number of task graphs in that data set

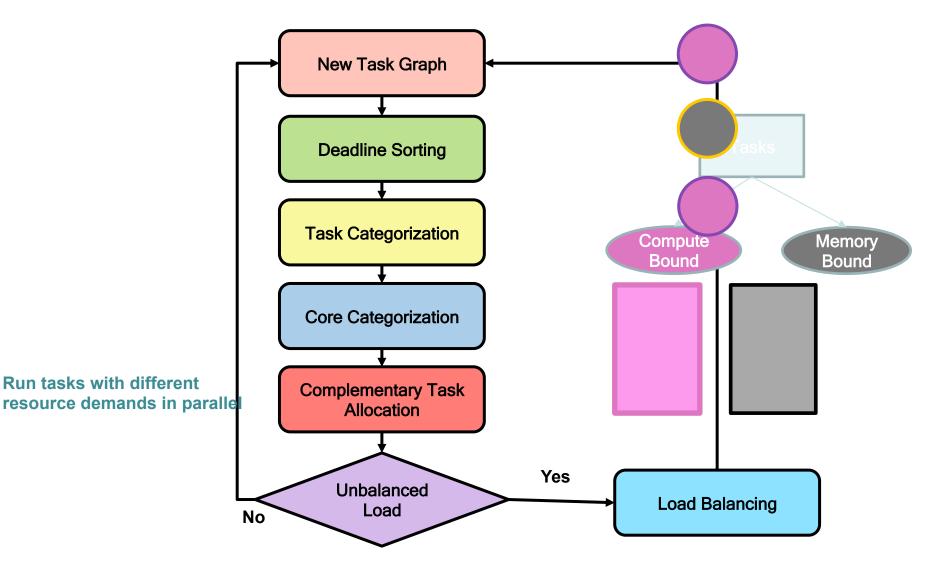
COMFAST



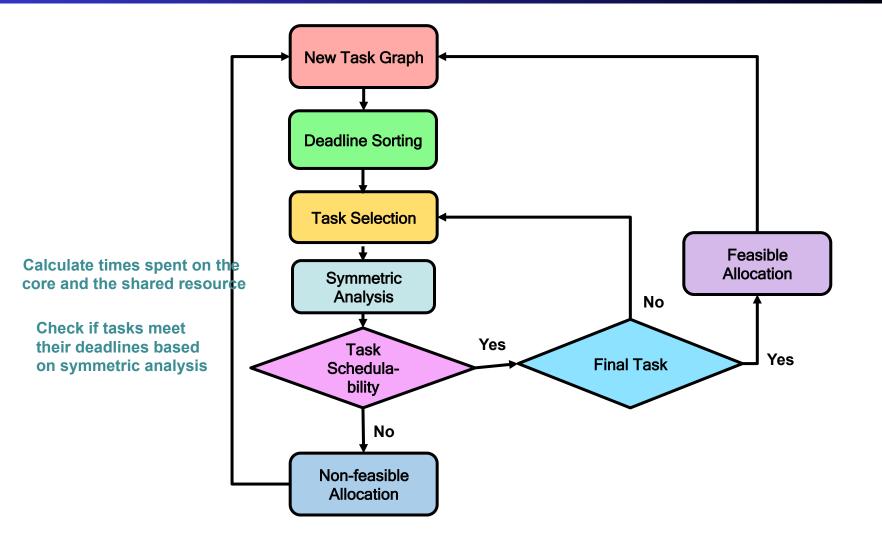
Case Studies

- Dynamic Task Aware Scheduling (DTAS) [Chiang, Journal of Systems and Software (JSS), 2016]
- First Fit (FF) Allocation Using Symmetric Analysis [Huang, Design Automation Conference (DAC), 2016]

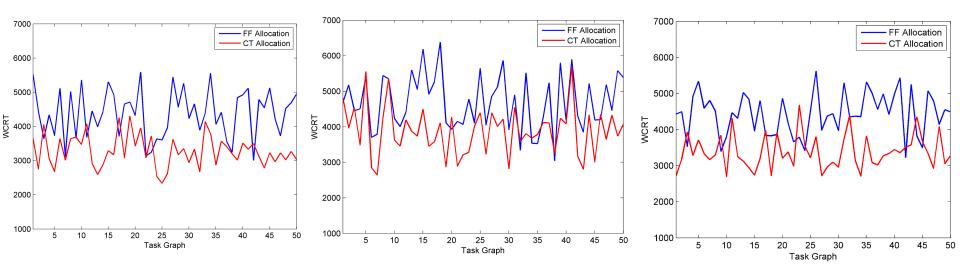
Dynamic Task Aware Scheduling (DTAS)



First Fit (FF) Allocation



Experimental Results: WCRT

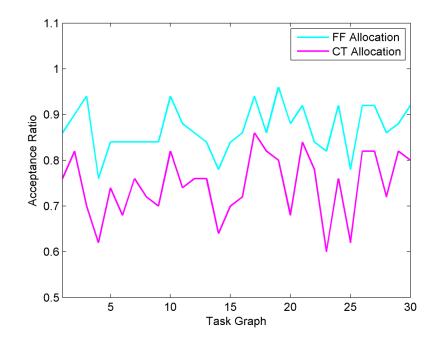


DTAS gives **better** results than FF Allocation

- 35.38% better for rare resource accesses
- 26.56% better for medium resource accesses
- 14.68% better for frequent resource accesses

Experimental Results: Acceptance Ratios

First Fit Allocation yields better results than complementary tasks by 15.3% at higher frequencies



COMFAST: A unified platform for comparison of scheduling techniques that focus primarily on minimizing contention

□ A generic data set

Task sets with varying resource requirements

Comparison is done based on

- Worst case response times
- Schedulability
- Ongoing and Future work
 - Optimization of scheduling algorithms based on our findings
 - □ More case studies

Thanks!

